

19. ANSWERS TO EXERCISES

Section 1.

(E1.1) Prove that if G is a finite simple abelian group, then $G \cong C_p$, the cyclic subgroup of order p , where p is a prime.

Answer. In an abelian group, any subgroup is normal. Thus a simple abelian group has no non-trivial proper subgroups. Suppose that G is non-cyclic and let $g \in G \setminus \{1\}$. Then, by supposition, $\langle g \rangle$ is proper and non-trivial and we have a contradiction. Suppose that $G = \langle g \rangle$ where $o(g) = st$ where s and t are positive integers greater than 1. Then $\langle g^s \rangle$ is a proper non-trivial subgroup of G and we have a contradiction. The result follows.

Section 2.

(E2.1) Prove that **Set**, **Pfn**, **Grp**, **Top** and **Vect_K** are categories.

Answer. Objects in each of these categories are structured sets. Let X be an object in the respective category, with Ω the underlying set. Since function composition is associative, the composition of arrows in these categories is associative and (C1) is satisfied. In addition the identity map $1 : \Omega \rightarrow \Omega$ ‘carries’ an arrow $X \rightarrow X$ that we call 1_X , and which satisfies (C2).

One should also check that the partial composition in **Pfn** is well-defined - one needs to admit the possibility that the function is defined on the empty set only!

(E2.2) Which categories in Example 1 are (full) subcategories of some other category in Example 1?

Answer.

- **Set** is a non-full subcategory of **Pfn**.
- **AGrp** is a full subcategory of **Grp**.
- **Field** is a full subcategory of **Ring**.
- **Vect_K** is a full subcategory of both **Mod – K** and **R – Mod**. (In fact these three categories are equivalent.)

(E2.3) Complete the definition of **Digraph** and prove that it is a category.

Answer. Objects: An object is a pair (V, E) where V is a set and E is a set of ordered pairs with entries from V . **Arrows:** An arrow

$$(V, E) \xrightarrow{f} (V', E')$$

is just a function $V \rightarrow V'$ such that

$$(e_1, e_2) \in E \implies (f(e_1), f(e_2)) \in E'.$$

Once again, since objects in **Digraph** are structured sets, (C1) and (C2) follow from the associativity of function composition, and the presence of an identity map on sets.

(E2.4) Give the ‘right’ definition of the category **Graph** corresponding to graphs that are not necessarily simple, i.e. which may have multiple edges between vertices.

Answer. Objects: An object is a triple (V, E, ι) where V and E are sets and $\iota : E \rightarrow 2^V$ is a function such that, for all $e \in E$, $\iota(e)$ has cardinality at most 2.

Arrows: An arrow

$$(V, E, \iota) \xrightarrow{f} (V', E', \iota')$$

is a pair of functions, $f_V : V \rightarrow V'$ and $f_E : E \rightarrow E'$, such that, for all $e \in E$,

$$\iota'(f_E(e)) = f_V(\iota(e)).$$

Note that the expression on the right hand side refers to the obvious induced map $f_V : 2^V \rightarrow 2^{V'}$.

(When we come to study isomorphisms we shall see why we cannot just extend the definition of **SimpleGraph** to this context.)

(E2.5) Prove that **Vect** \mathbf{S}_K and **IVect** \mathbb{R} are categories.

Answer. This is the same as previous answers for categories of structured sets.

(E2.6) Prove that **G-Set** is a category.

Answer. In this category, an arrow is a pair of functions. For a G -set (G, Ω, φ) we take the identity arrow to be $(1_G, 1_\Omega)$.

We need to check that the partial composition has the correct range of definition. Suppose we have two arrows as follows:

$$\begin{array}{ccc} G \times \Omega & \xrightarrow{\phi} & \Omega \\ (\alpha, \beta) \downarrow & & \downarrow \beta \\ H \times \Gamma & \xrightarrow{\psi} & \Gamma \end{array} \quad \begin{array}{ccc} H \times \Gamma & \xrightarrow{\psi} & \Gamma \\ (\gamma, \delta) \downarrow & & \downarrow \delta \\ J \times \Lambda & \xrightarrow{\xi} & \Lambda \end{array}$$

By definition, these two diagrams commute, hence if we consider the concatenated diagram

$$\begin{array}{ccc} G \times \Omega & \xrightarrow{\phi} & \Omega \\ (\alpha, \beta) \downarrow & & \downarrow \beta \\ H \times \Gamma & \xrightarrow{\psi} & \Gamma \\ (\gamma, \delta) \downarrow & & \downarrow \delta \\ J \times \Lambda & \xrightarrow{\xi} & \Lambda \end{array}$$

- then, since the two small rectangles commute, the large rectangle commutes. Now the pair $(\alpha\gamma, \beta\delta)$ is a well-defined arrow in **G-Set**, as required. Now (C1) and (C2) follow automatically.

(E2.7) Prove that Example 5 yields a category.

Answer. Clearly the partial composition is well-defined. For an object A , we define the identity arrow as follows:

$$1_A : A \times A \rightarrow \mathbb{R}, (a, b) \mapsto \begin{cases} 1, & \text{if } a = b; \\ 0, & \text{otherwise.} \end{cases}$$

To check (C1), suppose that the following arrows,

$$A \xrightarrow{f} B \xrightarrow{g} C \xrightarrow{h} D.$$

Now observe that

$$\begin{aligned} (fg)h : A \times D &\rightarrow \mathbb{R} \\ (a, d) &\mapsto \sum_{c \in C} fg(a, c)h(c, d) \\ &= \sum_{c \in C} \left(\sum_{b \in B} f(a, b)g(b, c) \right) h(c, d) \\ &= \sum_{b \in B} f(a, b) \left(\sum_{c \in C} g(b, c)h(c, d) \right) \\ &= \sum_{b \in B} f(a, b)gh(b, d). \end{aligned}$$

We conclude that $(fg)h = f(gh)$ as required. For (C2) consider the following arrows,

$$A \xrightarrow{1_A} A \xrightarrow{f} B \xrightarrow{1_B} B.$$

Now observe that

$$1_A f : A \times B \rightarrow \mathbb{R}, (a, b) \mapsto \sum_{a' \in A} 1_A(a, a')f(a', b) = f(a, b).$$

Thus $1_A f = f$ and, similarly, $f 1_B = f$ and we are done.

(E2.8) What is $\text{Aut}(X)$ when X is an object in **Top**?

Answer. The set of homeomorphisms of X .

(E2.9) Show that

- (1) an isomorphism is monic and epic;
- (2) if \mathbf{C} is a category of structured sets (so that each arrow is carried by a total function between the carriers of the two objects), then

injective \implies monic, and surjective \implies epic.

- (3) Show that epic does not imply surjective in **Ring**.
- (4) Show that bijective does not imply isomorphism in **Top**.

Answer.

- (1) Let $A \xrightarrow{f} B$ be an isomorphism and suppose that $X \xrightarrow{g} A$ are arrows such that $gf = hf$. Since f is an isomorphism, there exists $e : B \rightarrow A$ such that $fe = 1_B : B \rightarrow B$. Then

$$\begin{aligned} gf = hf &\implies (gf)e = (hf)e \\ &\implies g(fe) = h(fe) \\ &\implies g1_B = h1_B \\ &\implies g = h. \end{aligned}$$

We conclude that f is monic. The proof that f is epic is similar.